

# Quasi-Monte Carlo integration over a triangle

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# Integration over $\triangle$

- Let  $\triangle \subset \mathbb{R}^2$  be a triangle with Lebesgue measure denoted by  $|\triangle|$ .
- For an integrable  $f: \triangle \rightarrow \mathbb{R}$ , define

$$I(f) := \frac{1}{|\triangle|} \int_{\triangle} f(\mathbf{x}) \, d\mathbf{x}.$$

- We want to approximate  $I(f)$  using function evaluations.
- Such problem arises in graphics rendering and finite element methods.

# Quadrature over $\triangle$

- Consider a linear algorithm

$$A_N(f) = \sum_{n=0}^{N-1} w_n f(\mathbf{x}_n).$$

The Gauss-Legendre quadrature is popular for smooth  $f$ .

In this work we consider Quasi-Monte Carlo (QMC) integration

$$I_P(f) := \frac{1}{N} \sum_{n=0}^{N-1} f(\mathbf{x}_n),$$

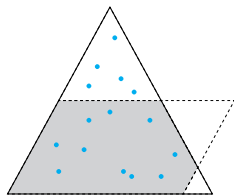
and give an explicit construction of  $P$  for twice differentiable  $f$ .

# Koksma-Hlawka

- Brandolini, Colzani, Gigante & Travaglini (2013) prove a Koksma-Hlawka type inequality for  $\triangle$  (more generally for simplices):

$$|I(f) - I_P(f)| \leq D^*(P)V(f).$$

- The discrepancy  $D^*(P)$  is the supremum of



among three vertices (set as origin).

# Review of Basu & Owen (2015)

- Our work is inspired by their one.
  - They come up with two constructions of low- $D^*(P)$  point sets in  $\Delta$ :
- ① Triangular van der Corput sequence:

$$D^*(P) = \frac{2}{3\sqrt{N}} - \frac{1}{9N}$$

for  $N = 4^\ell$ , extensible, and can be randomized.

- ② Triangular Kronecker lattices (rotation of an integer lattice through an angle whose tangent is badly approximable):

$$D^*(P) \leq C \frac{\log N}{N}.$$

## van der Corput (in base 4)

- Denote the 4-adic expansion of  $k \in \mathbb{N}_0$  by

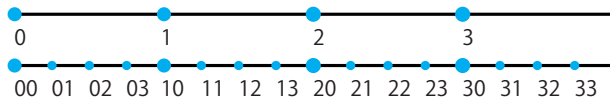
$$k = (\dots \kappa_2 \kappa_1 \kappa_0)_4$$

where all but a finite number of  $\kappa_j$ 's are 0.

- The original van der Corput sequence is given by  $\{\phi(k) : k \in \mathbb{N}_0\}$

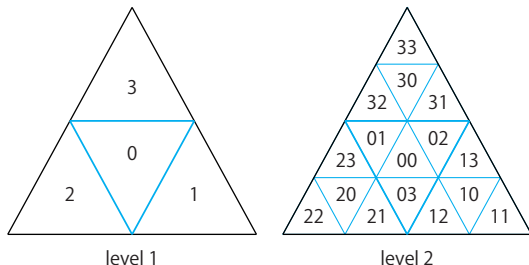
$$\phi(k) = (0.\kappa_0 \kappa_1 \kappa_2 \dots)_4$$

- This is based on a recursive partitioning of the unit interval:



# Triangular van der Corput

- Replace the recursive partitioning with the following one:



- $\kappa_0$  determines which subtriangle  $k$  is mapped to,  $\kappa_1$  does which sub-subtriangle  $k$  is mapped to, and so on.
- “Whether the subregion of level  $\ell$  is  $\triangle$  or  $\nabla$ ” is same as whether

$$\sigma_\ell(k) := \#\{0 \leq i \leq \ell - 1 : \kappa_i = 0\}$$

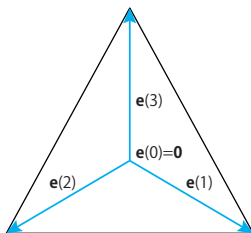
is even ( $\triangle$ ) or odd ( $\nabla$ ).

# Triangular van der Corput (cont'd)

- This idea gives the mapping

$$\tau(k) = \frac{\mathbf{e}(\kappa_0)}{2} + (-1)^{\sigma_1(k)} \frac{\mathbf{e}(\kappa_1)}{2^2} + (-1)^{\sigma_2(k)} \frac{\mathbf{e}(\kappa_2)}{2^3} + \dots$$

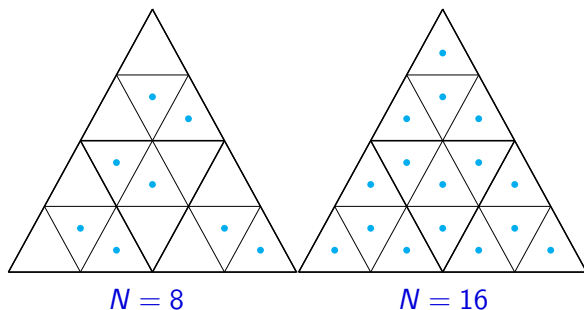
where



- The *triangular* van der Corput sequence is given by  $\{\tau(k) : k \in \mathbb{N}_0\}$ .



## Triangular van der Corput (cont'd)



- For  $N = 4^\ell$ , every subregion of level  $\ell$  has exactly one point in its center, giving

$$D^*(P) = \frac{2}{3\sqrt{N}} - \frac{1}{9N}$$

# Our idea

- ① Since we are dealing with 2D domains, it is natural to consider 2D QMC points or sequences instead of 1D van der Corput sequence.
- ② Focusing on smooth functions, there is a chance to conduct Walsh analysis as done in the unit cube setting (Dick, 2009; Yoshiki, 2017).
- ③ Given a decay of Walsh coefficients, it might be possible to prove the worst-case error bound.

# Class of functions and worst-case error

- In this work, we consider twice differentiable  $f$  with the norm

$$\|f\|_{C^2(\Delta)} := \max_{0 \leq \delta_1 + \delta_2 \leq 2} \left\| \frac{\partial^{\delta_1 + \delta_2} f}{\partial^{\delta_1} x_1 \partial^{\delta_2} x_2} \right\|_{L^\infty(\Delta)}.$$

- The worst-case error is defined by

$$e^{\text{wor}}(C^2(\Delta); A_N) = \sup_{\|f\|_{C^2(\Delta)} \leq 1} |I(f) - A_N(f)|.$$

For QMC integration using a point set  $P$ ,  $A_N$  is replaced by  $I_P$ .

## Digital sequences (in base 2)

- Let  $C_1, C_2 \in \mathbb{F}_2^{\mathbb{N} \times \mathbb{N}}$ , where  $\mathbb{F}_2$  denotes the two-element field.
- Denote the dyadic expansion of  $k \in \mathbb{N}_0$  by

$$k = (\dots \kappa_2 \kappa_1 \kappa_0)_2$$

where all but a finite number of  $\kappa_i$ 's are 0.

- Let

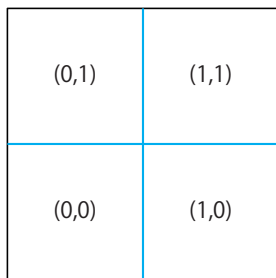
$$\begin{pmatrix} \xi_1^{(1)} \\ \xi_2^{(1)} \\ \vdots \end{pmatrix} = C_1 \begin{pmatrix} \kappa_0 \\ \kappa_1 \\ \vdots \end{pmatrix} \quad \text{and} \quad \begin{pmatrix} \xi_1^{(2)} \\ \xi_2^{(2)} \\ \vdots \end{pmatrix} = C_2 \begin{pmatrix} \kappa_0 \\ \kappa_1 \\ \vdots \end{pmatrix}.$$

- The digital sequence with  $C_1, C_2$  is given by

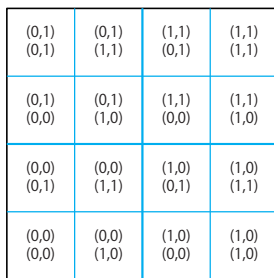
$$\left\{ ((0.\xi_1^{(1)}\xi_2^{(1)} \dots)_2, (0.\xi_1^{(2)}\xi_2^{(2)} \dots)_2) : k \in \mathbb{N}_0 \right\}.$$

# Recursive partitioning

- Digital sequences are based on a recursive partitioning of  $\square$ :



level 1



level 2

- $(\xi_1^{(1)}, \xi_1^{(2)})$  determines which subcube  $k$  is mapped to,  
 $(\xi_2^{(1)}, \xi_2^{(2)})$  does which sub-subcube  $k$  is mapped to, and so on.

## Quality criterion of $C_1, C_2$ for $\square$

- For simplicity, consider  $C_1, C_2 \in \mathbb{F}_2^{\ell \times \ell}$  for finite  $\ell$ . Define

$$P^\perp := \left\{ K = (\kappa_{ij}) \in \mathbb{F}_2^{\ell \times 2} : C_1^\top \begin{pmatrix} \kappa_{11} \\ \vdots \\ \kappa_{\ell 1} \end{pmatrix} \oplus C_2^\top \begin{pmatrix} \kappa_{12} \\ \vdots \\ \kappa_{\ell 2} \end{pmatrix} = \mathbf{0} \right\}.$$

- For  $k = (\kappa_i) \in \mathbb{F}_2^\ell$ , define

$$\mu_{\square}(k) = \begin{cases} 0 & \text{if } k = \mathbf{0}, \\ \max\{i \leq \ell : \kappa_i \neq 0\} & \text{otherwise.} \end{cases}$$

For  $K = (k_1, k_2) \in \mathbb{F}_2^{\ell \times 2}$ ,  $\mu_{\square}(K) := \mu_{\square}(k_1) + \mu_{\square}(k_2)$ .

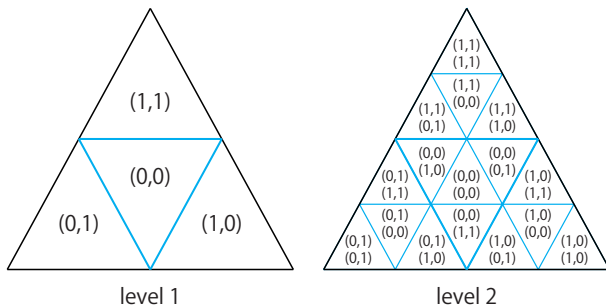
- So-called  $t$ -value is defined as the minimum  $t$  such that

$$\min_{K \in P^\perp \setminus \{\mathbf{0}\}} \mu_{\square}(K) \geq \ell - t + 1.$$

- Construction of  $C_1, C_2$  with small  $t$ -value: Sobol' (1967), Faure (1982), Niederreiter (1988), etc.

# Triangular digital sequences

- Replace the recursive partitioning with the following one:



- “Whether the subregion of level  $\ell$  is  $\triangle$  or  $\nabla$ ” is same as whether

$$\sigma_\ell(k) := \#\{1 \leq i \leq \ell : (\xi_i^{(1)}, \xi_i^{(2)}) = (0, 0)\}$$

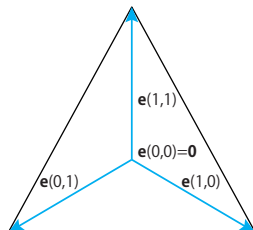
is even ( $\triangle$ ) or odd ( $\nabla$ ).

## Triangular digital sequences (cont'd)

- The *triangular* digital sequence with  $C_1, C_2$  is  $\{\tau(k) : k \in \mathbb{N}_0\}$  with

$$\tau(k) = \frac{\mathbf{e}(\xi_1^{(1)}, \xi_1^{(2)})}{2} + (-1)^{\sigma_1(k)} \frac{\mathbf{e}(\xi_2^{(1)}, \xi_2^{(2)})}{2^2} + (-1)^{\sigma_2(k)} \frac{\mathbf{e}(\xi_3^{(1)}, \xi_3^{(2)})}{2^3} + \dots$$

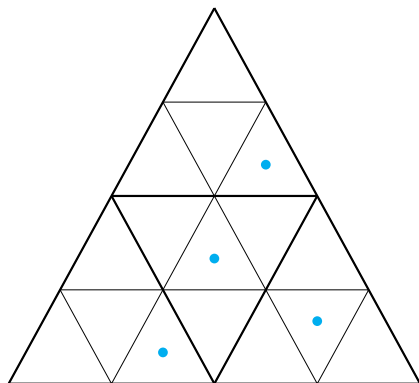
where



- The triangular regular grids is same as the triangular van der Corput.  
→ Our construction includes Basu-Owen's one as a special case.

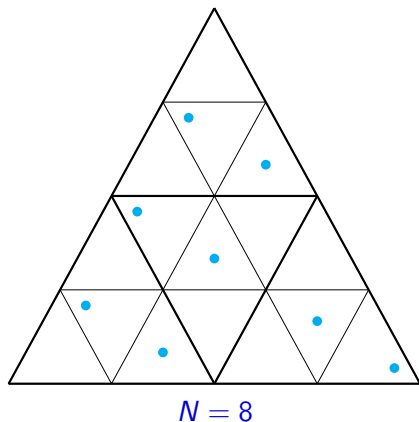


# Triangular Sobol' sequences

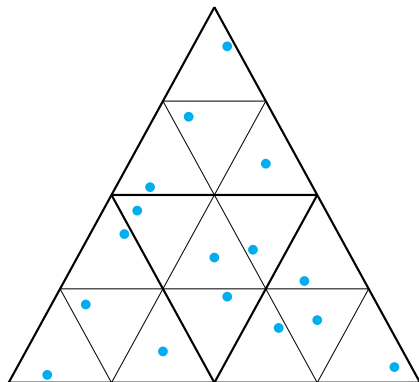


$N = 4$

# Triangular Sobol' sequences

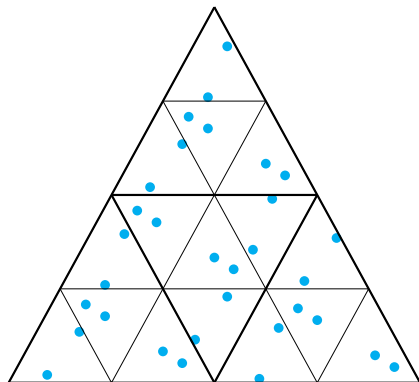


# Triangular Sobol' sequences



$N = 16$

# Triangular Sobol' sequences



$$N = 32$$

# Discretization

- For a fixed level  $\ell$ , the set of  $4^\ell$  subregions and  $\mathbb{F}_2^{\ell \times 2}$  are one-to-one.
- Denote the subregion corresponding to  $X \in \mathbb{F}_2^{\ell \times 2}$  by  $\Delta_\ell(X)$ . Define

$$f_\ell(X) = \frac{1}{4^\ell |\Delta|} \int_{\Delta_\ell(X)} f(\mathbf{x}) \, d\mathbf{x}.$$

## Lemma

For  $X \in \mathbb{F}_2^{\ell \times 2}$  and  $\mathbf{x} \in \Delta_\ell(X)$ , we have

$$|f(\mathbf{x}) - f_\ell(X)| \leq \frac{\sqrt{2}d(\Delta)\|f\|_{C^2(\Delta)}}{2^\ell}$$

where  $d(\Delta)$  denotes the diameter of  $\Delta$ .

# Walsh series

- For  $X = (\xi_{ij})$ ,  $K = (\kappa_{ij}) \in \mathbb{F}_2^{\ell \times 2}$ , define

$$\text{wal}_K(X) := (-1)^{\sum_{i,j} \kappa_{ij} \xi_{ij}} = \prod_{i,j} (-1)^{\kappa_{ij} \xi_{ij}},$$

called the  $K$ -th Walsh function.

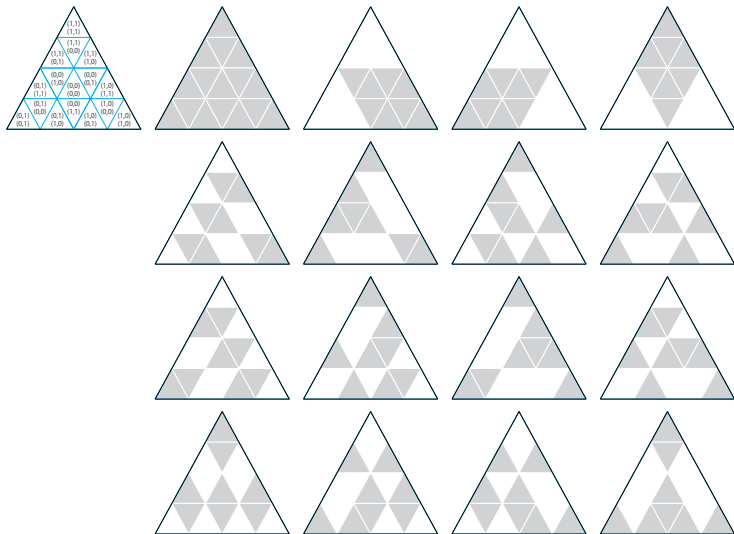
- For  $f_\ell: \mathbb{F}_2^{\ell \times 2} \rightarrow \mathbb{R}$ , we have a Walsh series expansion:

$$f_\ell(X) = \sum_{K \in \mathbb{F}_2^{\ell \times 2}} \hat{f}_\ell(K) \text{wal}_K(X)$$

where

$$\hat{f}_\ell(K) = \frac{1}{4^\ell} \sum_{X \in \mathbb{F}_2^{\ell \times 2}} f_\ell(X) \text{wal}_K(X).$$

# Triangular Walsh functions



# Decay of Walsh coefficients

- For  $K = (k_1, k_2) \in \mathbb{F}_2^{\ell \times 2}$ , define

$$\mu_{\Delta}(K) = \max \{ \mu_{\square}(k_1), \mu_{\square}(k_2) \}.$$

(Recall:  $\mu_{\square}(K) := \mu_{\square}(k_1) + \mu_{\square}(k_2)$ )

## Lemma

For  $\ell \in \mathbb{N}$  and  $K \in \mathbb{F}_2^{\ell \times 2}$ , we have

$$|\hat{f}_{\ell}(K)| \leq \max \left( 2\sqrt{2}d(\Delta), 4d^2(\Delta) \right) \|f\|_{C^2(\Delta)} \frac{\mu_{\Delta}(K)}{2^{\mu_{\Delta}(K)}}.$$



## Quality criterion of $C_1, C_2$ for $\Delta$

- $C_1, C_2$  with small  $t$ -value is already OK because

$$\min_{K \in P^\perp \setminus \{0\}} \mu_\Delta(K) \geq \min_{K \in P^\perp \setminus \{0\}} \frac{\mu_\square(K)}{2} \geq \frac{\ell - t + 1}{2}$$

- The triangular regular grids is also good in this regard. Thus the triangular van der Corput sequence due to Basu & Owen (2015) is not optimal in terms of  $D^*(P)$ , but almost optimal for numerical integration of  $f \in C^2(\Delta)$  (as shown below).

# Upper bound

- The following is the main result of this work:

Theorem (G., Suzuki and Yoshiki, 2017)

*QMC integration using a triangular digital sequence with large minimum  $\mu_\Delta$ -value achieves*

$$e^{\text{wor}}(C^2(\Delta); I_P) \leq C \frac{(\log N)^3}{N}.$$

*Moreover, for  $N = 2^\ell$  with  $\ell \in \mathbb{N}$ , we have*

$$e^{\text{wor}}(C^2(\Delta); I_P) \leq C \frac{\ell^2}{2^\ell}.$$

## Lower bound

- The bump function technique from Bakhvalov (1959) gives:

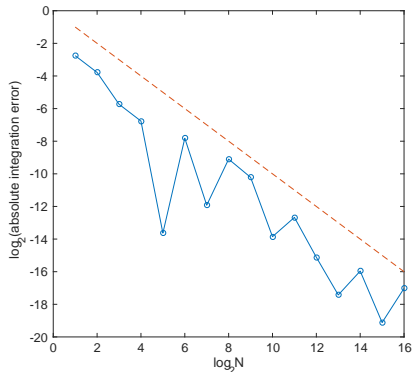
### Theorem

For any  $A_N$ , there exists a constant  $c > 0$  s.t.

$$e^{\text{wor}}(C^2(\Delta); A_N) \geq \frac{c}{N}.$$

# Numerics

- Consider  $\Delta = \{(x, y) : x, y \geq 0, x + y \leq 1\}$  and use triangular Sobol' sequences.
- The test function is  $f(x, y) = e^x$  for which we have  $I(f) = 2(e - 2)$ .
- The red line shows  $N^{-1}$  as a reference.



Thank you for your attention!